Segment Routing Configuration Guide for Cisco NCS 5500 Series Routers, IOS XR 6.2.x

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Preface

Note This product has reached end-of-life status. For more information, see the End-of-Life and End-of-Sale Notices.

The *Segment Routing Configuration Guide for for Cisco NCS 5500 Series Routers* preface contains these sections:

Changes to This Document

This table lists the changes made to this document since it was first printed.

Date	**Change Summary**
March 2017	Initial release of this document

Communications, Services, and Additional Information

- To receive timely, relevant information from Cisco, sign up at Cisco Profile Manager.
- To get the business impact you're looking for with the technologies that matter, visit Cisco Services.
- To submit a service request, visit Cisco Support.
- To discover and browse secure, validated enterprise-class apps, products, solutions and services, visit Cisco Marketplace.
- To obtain general networking, training, and certification titles, visit Cisco Press.
- To find warranty information for a specific product or product family, access Cisco Warranty Finder.

Cisco Bug Search Tool

Cisco Bug Search Tool (BST) is a web-based tool that acts as a gateway to the Cisco bug tracking system that maintains a comprehensive list of defects and vulnerabilities in Cisco products and software. BST provides you with detailed defect information about your products and software.

CHAPTER 1

New and Changed Information for Segment Routing Features

This table summarizes the new and changed feature information for the *Segment Routing Configuration Guide for Cisco NCS 5500 Series Routers*, and lists where they are documented.

- New and Changed Information, on page 1

New and Changed Information

Table 1: New and Changed Features

Feature	**Description**	**Introduced/Changed in Release**	**Where Documented**
IS-IS Multi-Domain Prefix SID and Domain Stitching	This feature was introduced.	Release 6.2.2	*Configure Segment Routing for IS-IS Protocol*
IOS XR Traffic Controller (XTC)	This feature was introduced.	Release 6.2.2	*Configure IOS XR Traffic Controller (XTC)*
MPLS Ping and Traceroute for Prefix-SID	This feature was introduced.	Release 6.2.2	*Using Segment Routing OAM*

CHAPTER 2

About Segment Routing

This chapter introduces the concept of segment routing and provides a workflow for configuring segment routing.

Scope

Segment routing is a method of forwarding packets on the network based on the source routing paradigm. The source chooses a path and encodes it in the packet header as an ordered list of segments. Segments are an identifier for any type of instruction. For example, topology segments identify the next hop toward a destination. Each segment is identified by the segment ID (SID) consisting of a flat unsigned 20-bit integer.

Segments

Interior gateway protocol (IGP) distributes two types of segments: prefix segments and adjacency segments. Each router (node) and each link (adjacency) has an associated segment identifier (SID).

- A prefix SID is associated with an IP prefix. The prefix SID is manually configured from the segment routing global block (SRGB) range of labels, and is distributed by IS-IS or OSPF. The prefix segment steers the traffic along the shortest path to its destination. A node SID is a special type of prefix SID that identifies a specific node. It is configured under the loopback interface with the loopback address of the node as the prefix.

  A prefix segment is a global segment, so a prefix SID is globally unique within the segment routing domain.

- An adjacency segment is identified by a label called an adjacency SID, which represents a specific adjacency, such as egress interface, to a neighboring router. The adjacency SID is distributed by IS-IS or OSPF. The adjacency segment steers the traffic to a specific adjacency.

  An adjacency segment is a local segment, so the adjacency SID is locally unique relative to a specific router.

By combining prefix (node) and adjacency segment IDs in an ordered list, any path within a network can be constructed. At each hop, the top segment is used to identify the next hop. Segments are stacked in order at the top of the packet header. When the top segment contains the identity of another node, the receiving node uses equal cost multipaths (ECMP) to move the packet to the next hop. When the identity is that of the receiving node, the node pops the top segment and performs the task required by the next segment.

Dataplane

Segment routing can be directly applied to the Multiprotocol Label Switching (MPLS) architecture with no change in the forwarding plane. A segment is encoded as an MPLS label. An ordered list of segments is encoded as a stack of labels. The segment to process is on the top of the stack. The related label is popped from the stack, after the completion of a segment.

Services

Segment Routing integrates with the rich multi-service capabilities of MPLS, including Layer 3 VPN (L3VPN), Virtual Private Wire Service (VPWS), Virtual Private LAN Service (VPLS), and Ethernet VPN (EVPN).

Segment Routing for Traffic Engineering

Segment routing for traffic engineering (SR-TE) takes place through a tunnel between a source and destination pair. Segment routing for traffic engineering uses the concept of source routing, where the source calculates the path and encodes it in the packet header as a segment. Each segment is an end-to-end path from the source to the destination, and instructs the routers in the provider core network to follow the specified path instead of the shortest path calculated by the IGP. The destination is unaware of the presence of the tunnel.

Need

With segment routing for traffic engineering (SR-TE), the network no longer needs to maintain a per-application and per-flow state. Instead, it simply obeys the forwarding instructions provided in the packet.

SR-TE utilizes network bandwidth more effectively than traditional MPLS-TE networks by using ECMP at every segment level. It uses a single intelligent source and relieves remaining routers from the task of calculating the required path through the network.

Benefits

- **Ready for SDN**: Segment routing was built for SDN and is the foundation for Application Engineered Routing (AER). SR prepares networks for business models, where applications can direct network behavior. SR provides the right balance between distributed intelligence and centralized optimization and programming.
- **Minimal configuration**: Segment routing for TE requires minimal configuration on the source router.
- **Load balancing**: Unlike in RSVP-TE, load balancing for segment routing can take place in the presence of equal cost multiple paths (ECMPs).
- **Supports Fast Reroute (FRR)**: Fast reroute enables the activation of a pre-configured backup path within 50 milliseconds of path failure.
- **Plug-and-Play deployment**: Segment routing tunnels are interoperable with existing MPLS control and data planes and can be implemented in an existing deployment.

Workflow for Deploying Segment Routing

Follow this workflow to deploy segment routing.

1. Configure the Segment Routing Global Block (SRGB)
2. Enable Segment Routing and Node SID for the IGP
3. Configure Segment Routing for BGP
4. Configure the SR-TE Policy
5. Configure the SR-PCE
6. Configure TI-LFA and Microloop Avoidance
7. Configure the Segment Routing Mapping Server

CHAPTER 3

Configure Segment Routing Global Block

Local label allocation is managed by the label switching database (LSD). The Segment Routing Global Block (SRGB) is the range of label values preserved for segment routing in the LSD.

- About the Segment Routing Global Block, on page 7
- Setup a Non-Default Segment Routing Global Block Range, on page 8

About the Segment Routing Global Block

The SRGB label values are assigned as prefix segment identifiers (SIDs) to SR-enabled nodes and have global significance throughout the domain.

Note Because the values assigned from the range have domain-wide significance, we recommend that all routers within the domain be configured with the same range of values.

On SR-capable routers, the default starting value of the dynamic label range is changed from 16000 to 24000, so that the default SRGB label values (16000 to 23999) are available when SR is enabled on a running system. If a dynamic label range has been configured with a starting value of 16000, then the default SRGB label values may already be in use when SR is enabled on a running system. Therefore, you must reload the router after enabling SR to release the currently allocated labels and allocate the SRGB.

Also, if you need to increase the SRGB range after you have enabled SR, you must reload the router to release the currently allocated labels and allocate the new SRGB.

To keep the segment routing configuration simple and to make it easier to troubleshoot segment routing issues, we recommend that you use the default SRGB range on each node in the domain. However, there are instances when you might need to define a different range. For example:

- The nodes of another vendor support a label range that is different from the default SRGB, and you want to use the same SRGB on all nodes.
- The default range is too small.
- To specify separate SRGBs for IS-IS and OSPF protocols, as long as the ranges do not overlap.

The SRGB can be disabled if SR is not used.

Behaviors and Limitations

- The default SRGB has a size of 8000 starting from label value 16000; therefore, the default SRGB range goes from 16000 to 23,999.

Note Label values that are not previously reserved are available for dynamic assignment.

Setup a Non-Default Segment Routing Global Block Range

This task explains how to configure a non-default SRGB range.

SUMMARY STEPS

1. **configure**
2. [**router** {**isis** *instance-id* | **ospf** *process_name*}]
3. **segment-routing global-block** *starting_value ending_value*
4. Use the **commit** or **end** command.

DETAILED STEPS

	Command or Action	Purpose
Step 1	**configure** **Example:** `RP/0/RP0/CPU0:router# configure`	Enters mode.
Step 2	[**router** {**isis** *instance-id* | **ospf** *process_name*}] **Example:** `RP/0/RP0/CPU0:router(config)# router isis 1`	(Optional) Enter the **router isis** *instance-id* or **router ospf** *process_name* commands if you want to configure separate SRGBs for IS-IS and OSPF protocols.
Step 3	**segment-routing global-block** *starting_value ending_value* **Example:** `RP/0/RP0/CPU0:router(config-isis)# segment-routing global-block 18000 19999`	Enter the lowest value that you want the SRGB range to include as the starting value. Enter the highest value that you want the SRGB range to include as the ending value.
Step 4	Use the **commit** or **end** command.	**commit** —Saves the configuration changes and remains within the configuration session. **end** —Prompts user to take one of these actions: • **Yes** — Saves configuration changes and exits the configuration session.

	Command or Action	Purpose
		- **No** —Exits the configuration session without committing the configuration changes. - **Cancel** —Remains in the configuration session, without committing the configuration changes.

Verify the SRGB configuration:

```
RP/0/RP0/CPU0:router# show mpls label table detail
Table Label   Owner                           State  Rewrite
----- ------- ------------------------------- ------ -------

<...snip...>

0     18000   ISIS(A):1                       InUse  No
   Lbl-blk SRGB, vers:0, (start_label=18000, size=2000)
0     24000   ISIS(A):1                       InUse  Yes
   (SR Adj Segment IPv4, vers:0, index=1, type=0, intf=Gi0/0/0/0, nh=10.0.0.2)
```

What to do next

Configure prefix SIDs and enable segment routing.

CHAPTER 4

Configure Segment Routing for IS-IS Protocol

Integrated Intermediate System-to-Intermediate System (IS-IS), Internet Protocol Version 4 (IPv4), is a standards-based Interior Gateway Protocol (IGP). The Cisco IOS XR software implements the IP routing capabilities described in International Organization for Standardization (ISO)/International Engineering Consortium (IEC) 10589 and RFC 1995, and adds the standard extensions for single topology and multitopology IS-IS for IP Version 6 (IPv6).

This module provides the configuration information used to enable segment routing for IS-IS.

Note For additional information on implementing IS-IS on your Cisco NCS 5500 Series Router, see the *Implementing IS-IS* module in the *Routing Configuration Guide for Cisco NCS 5500 Series Routers*.

Enabling Segment Routing for IS-IS Protocol

Segment routing on the IS-IS control plane supports the following:

- IPv4 and IPv6 control plane
- Level 1, level 2, and multi-level routing
- Prefix SIDs for host prefixes on loopback interfaces
- Adjacency SIDs for adjacencies
- MPLS penultimate hop popping (PHP) and explicit-null signaling

This task explains how to enable segment routing for IS-IS.

Before you begin

Your network must support the MPLS Cisco IOS XR software feature before you enable segment routing for IS-IS on your router.

Note You must enter the commands in the following task list on every IS-IS router in the traffic-engineered portion of your network.

SUMMARY STEPS

1. **configure**
2. **router isis** *instance-id*
3. **address-family** { **ipv4** | **ipv6** } [**unicast**]
4. **metric-style wide** [**level** { **1** | **2** }]
5. **mpls traffic-eng** *level*
6. **mpls traffic-eng router-id** *interface*
7. **router-id loopback** *loopback interface used for prefix-sid*
8. **segment-routing mpls**
9. **exit**
10. **mpls traffic-eng**
11. Use the **commit** or **end** command.

DETAILED STEPS

	Command or Action	Purpose
Step 1	**configure** **Example:** `RP/0/RP0/CPU0:router# configure`	Enters mode.
Step 2	**router isis** *instance-id* **Example:** `RP/0/RP0/CPU0:router(config)# router isis isp`	Enables IS-IS routing for the specified routing instance, and places the router in router configuration mode. **Note** You can change the level of routing to be performed by a particular routing instance by using the **is-type** router configuration command.
Step 3	**address-family** { **ipv4** | **ipv6** } [**unicast**] **Example:** `RP/0/RP0/CPU0:router(config-isis)# address-family ipv4 unicast`	Specifies the IPv4 or IPv6 address family, and enters router address family configuration mode.
Step 4	**metric-style wide** [**level** { **1** | **2** }] **Example:** `RP/0/RP0/CPU0:router(config-isis-af)# metric-style wide level 1`	Configures a router to generate and accept only wide link metrics in the Level 1 area.
Step 5	**mpls traffic-eng** *level* **Example:**	Enables RSVP traffic engineering funtionality.

	Command or Action	Purpose
	`RP/0/RP0/CPU0:router(config-isis-af)# mpls traffic-eng level-2-only`	
Step 6	**mpls traffic-eng router-id** *interface* **Example:** `RP/0/RP0/CPU0:router(config-isis-af)# mpls traffic-eng router-id Loopback0`	Sets the traffic engineering loopback interface.
Step 7	**router-id loopback** *loopback interface used for prefix-sid* **Example:** `RP/0/RP0(config-isis-af)#router-id loopback0`	Configures router ID for each address-family (ipv4/ipv6).
Step 8	**segment-routing mpls** **Example:** `RP/0/RP0/CPU0:router(config-isis-af)# segment-routing mpls`	Segment routing is enabled by the following actions: • MPLS forwarding is enabled on all interfaces where IS-IS is active. • All known prefix-SIDs in the forwarding plain are programmed, with the prefix-SIDs advertised by remote routers or learned through local or remote mapping server. • The prefix-SIDs locally configured are advertised.

Step 9	**exit**
	Example:
	`RP/0/RP0/CPU0:router(config-isis-af)# exit`
	`RP/0/RP0/CPU0:router(config-isis)# exit`	

Step 10	**mpls traffic-eng** **Example:** `RP/0/RP0/CPU0:router(config)# mpls traffic-eng`	Enables traffic engineering functionality on the node. The node advertises the traffic engineering link attributes in IGP which populates the traffic engineering database (TED) on the head-end. The RSVP-TE head-end requires the TED to calculate and validate the path of the RSVP-TE policy.
Step 11	Use the **commit** or **end** command.	**commit** —Saves the configuration changes and remains within the configuration session. **end** —Prompts user to take one of these actions: • **Yes** — Saves configuration changes and exits the configuration session. • **No** —Exits the configuration session without committing the configuration changes. • **Cancel** —Remains in the configuration session, without committing the configuration changes.

What to do next

Configure the prefix SID.

Configuring a Prefix-SID on the IS-IS Enabled Loopback Interface

A prefix segment identifier (SID) is associated with an IP prefix. The prefix SID is manually configured from the segment routing global block (SRGB) range of labels. A prefix SID is configured under the loopback interface with the loopback address of the node as the prefix. The prefix segment steers the traffic along the shortest path to its destination.

A prefix SID can be a node SID or an Anycast SID. A node SID is a type of prefix SID that identifies a specific node. An Anycast SID is a type of prefix SID that identifies a set of nodes, and is configured with n-flag clear. The set of nodes (Anycast group) is configured to advertise a shared prefix address and prefix SID. Anycast routing enables the steering of traffic toward multiple advertising nodes. Packets addressed to an Anycast address are forwarded to the topologically nearest nodes.

The prefix SID is globally unique within the segment routing domain.

This task explains how to configure prefix segment identifier (SID) index or absolute value on the IS-IS enabled Loopback interface.

Before you begin

Ensure that segment routing is enabled on the corresponding address family.

SUMMARY STEPS

1. **configure**
2. **router isis** *instance-id*
3. **interface Loopback** *instance*
4. **address-family** { **ipv4** | **ipv6** } [**unicast**]
5. **prefix-sid** {**index** *SID-index* | **absolute** *SID-value*} [**n-flag-clear**] [**explicit-null**]
6. Use the **commit** or **end** command.

DETAILED STEPS

	Command or Action	Purpose
Step 1	**configure** **Example:** `RP/0/RP0/CPU0:router# configure`	Enters mode.
Step 2	**router isis** *instance-id* **Example:** `RP/0/RP0/CPU0:router(config)# router isis 1`	Enables IS-IS routing for the specified routing instance, and places the router in router configuration mode. • You can change the level of routing to be performed by a particular routing instance by using the **is-type** router configuration command.

	Command or Action	Purpose
Step 3	**interface Loopback** *instance* **Example:** `RP/0/RP0/CPU0:router(config-isis)# interface Loopback0`	Specifies the loopback interface and instance.
Step 4	**address-family** { **ipv4** | **ipv6** } [**unicast**] **Example:** The following is an example for ipv4 address family: `RP/0/RP0/CPU0:router(config-isis-if)# address-family ipv4 unicast`	Specifies the IPv4 or IPv6 address family, and enters router address family configuration mode.

Step 5	**prefix-sid** {**index** *SID-index* | **absolute** *SID-value*} [**n-flag-clear**] [**explicit-null**]
	Example:
	`RP/0/RP0/CPU0:router(config-isis-if-af)# prefix-sid index 1001`
	`RP/0/RP0/CPU0:router(config-isis-if-af)# prefix-sid absolute 17001`
		Configures the prefix-SID index or absolute value for the interface.
		Specify **index** *SID-index* for each node to create a prefix SID based on the lower boundary of the SRGB + the index.
		Specify **absolute** *SID-value* for each node to create a specific prefix SID within the SRGB.
		By default, the n-flag is set on the prefix-SID, indicating that it is a node SID. For specific prefix-SID (for example, Anycast prefix-SID), enter the `n-flag-clear` keyword. IS-IS does not set the `N` flag in the prefix-SID sub Type Length Value (TLV).
		To disable penultimate-hop-popping (PHP) and add explicit-Null label, enter `explicit-null` keyword. IS-IS sets the `E` flag in the prefix-SID sub TLV.

Step 6	Use the **commit** or **end** command.	**commit** —Saves the configuration changes and remains within the configuration session. **end** —Prompts user to take one of these actions: • **Yes** — Saves configuration changes and exits the configuration session. • **No** —Exits the configuration session without committing the configuration changes. • **Cancel** —Remains in the configuration session, without committing the configuration changes.

Verify the prefix-SID configuration:

```
RP/0/RP0/CPU0:router# show isis database verbose

IS-IS 1 (Level-2) Link State Database
LSPID                 LSP Seq Num  LSP Checksum  LSP Holdtime  ATT/P/OL
router.00-00        * 0x0000039b   0xfc27        1079            0/0/0
```

```
  Area Address: 49.0001
  NLPID:        0xcc
  NLPID:        0x8e
  MT:           Standard (IPv4 Unicast)
  MT:           IPv6 Unicast                                 0/0/0
  Hostname:     router
  IP Address:   10.0.0.1
  IPv6 Address: 2001:0db8:1234::0a00:0001
  Router Cap:   10.0.0.1, D:0, S:0
    Segment Routing: I:1 V:1, SRGB Base: 16000 Range: 8000
    SR Algorithm:
      Algorithm: 0

<...>
  Metric: 0           IP-Extended 10.0.0.1/32
    Prefix-SID Index: 1001, Algorithm:0, R:0 N:1 P:0 E:0 V:0 L:0

<...>
```

What to do next

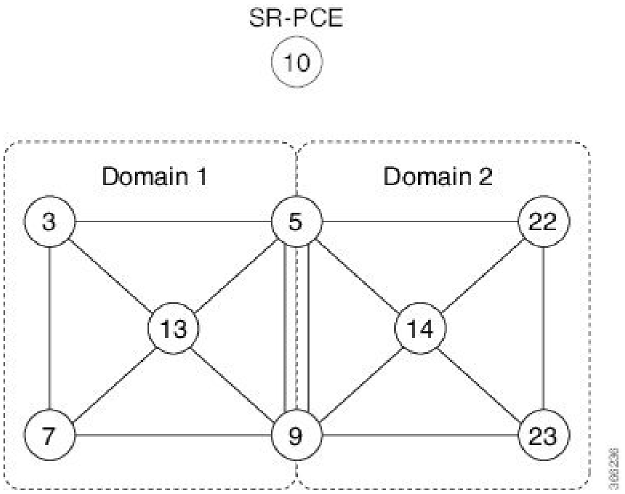
Configure the SR-TE policy.

IS-IS Multi-Domain Prefix SID and Domain Stitching: Example

IS-IS Multi-Domain Prefix SID and Domain Stitching allows you to configure multiple IS-IS instances on the same loopback interface for domain border nodes. You specify a loopback interface and prefix SID under multiple IS-IS instances to make the prefix and prefix SID reachable in different domains.

This example uses the following topology. Node 5 and 9 are border nodes between two IS-IS domains (Domain1 and Domain2). Node 10 is configured as the Segment Routing Path Computation Element (SR-PCE) (see Configure Segment Routing Path Computation Element).

Figure 1: Multi-Domain Topology



Configure IS-IS Multi-Domain Prefix SID

Specify a loopback interface and prefix SID under multiple IS-IS instances on each border node:

Example: Border Node 5

```
router isis Domain1
 interface Loopback0
  address-family ipv4 unicast
   prefix-sid absolute 16005

router isis Domain2
 interface Loopback0
  address-family ipv4 unicast
   prefix-sid absolute 16005
```

Example: Border Node 9

```
router isis Domain1
 interface Loopback0
  address-family ipv4 unicast
   prefix-sid absolute 16009

router isis Domain2
 interface Loopback0
  address-family ipv4 unicast
   prefix-sid absolute 16009
```

Border nodes 5 and 9 each run two IS-IS instances (Domain1 and Domain2) and advertise their Loopback0 prefix and prefix SID in both domains.

Nodes in both domains can reach the border nodes by using the same prefix and prefix SID. For example, Node 3 and Node 22 can reach Node 5 using prefix SID 16005.

Configure Common Router ID

On each border node, configure a common TE router ID under each IS-IS instance:

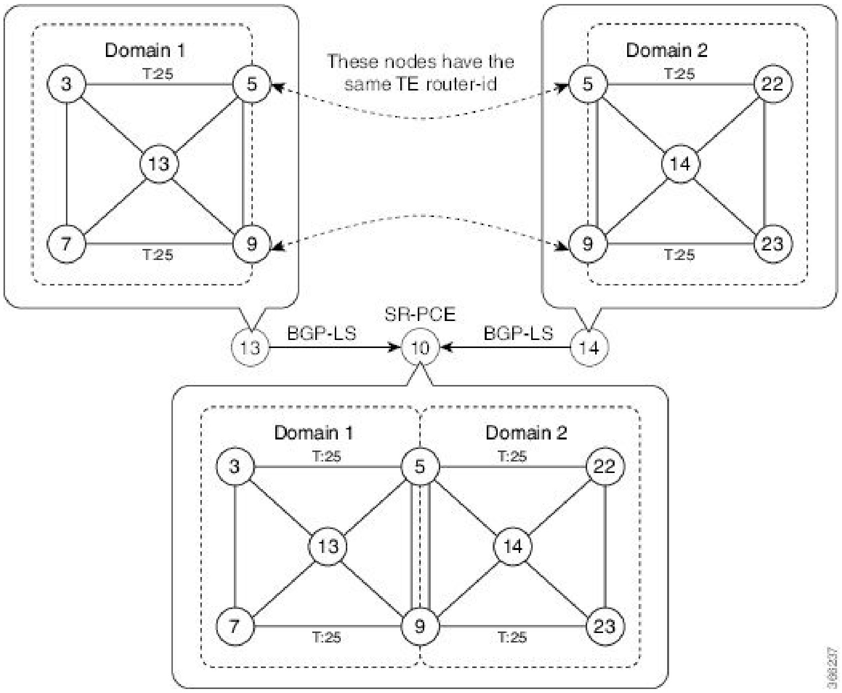
Example: Border Node 5

```
router isis Domain1
 address-family ipv4 unicast
  router-id loopback0

router isis Domain2
 address-family ipv4 unicast
  router-id loopback0
```

Example: Border Node 9

```
router isis Domain1
 address-family ipv4 unicast
  router-id loopback0

router isis Domain2
 address-family ipv4 unicast
  router-id loopback0
```

Distribute IS-IS Link-State Data



Configure BGP Link-state (BGP-LS) on Node 13 and Node 14 to report their local domain to Node 10:

```
Example: Node 13
router isis Domain1
 distribute link-state instance-id instance-id
```

```
Example: Node 14
router isis Domain2
 distribute link-state instance-id instance-id
```

Link-state ID starts from 32. One ID is required per IGP domain. Different domain IDs are essential to identify that the SR-TE TED belongs to a particular IGP domain.

Nodes 13 and 14 each reports its local domain in BGP-LS to Node 10.

Node 10 identifies the border nodes (Nodes 5 and 9) by their common advertised TE router ID, then combines (stitches) the domains on these border nodes for end-to-end path computations.

CHAPTER 5

Configure Segment Routing for OSPF Protocol

Open Shortest Path First (OSPF) is an Interior Gateway Protocol (IGP) developed by the OSPF working group of the Internet Engineering Task Force (IETF). Designed expressly for IP networks, OSPF supports IP subnetting and tagging of externally derived routing information. OSPF also allows packet authentication and uses IP multicast when sending and receiving packets.

This module provides the configuration information to enable segment routing for OSPF.

Note For additional information on implementing OSPF on your , see the *Implementing OSPF* module in the .

Enabling Segment Routing for OSPF Protocol

Segment routing on the OSPF control plane supports the following:

- OSPFv2 control plane
- Multi-area
- IPv4 prefix SIDs for host prefixes on loopback interfaces
- Adjacency SIDs for adjacencies
- MPLS penultimate hop popping (PHP) and explicit-null signaling

This section describes how to enable segment routing MPLS and MPLS forwarding in OSPF. Segment routing can be configured at the instance, area, or interface level.

Before you begin

Your network must support the MPLS Cisco IOS XR software feature before you enable segment routing for OSPF on your router.

Note You must enter the commands in the following task list on every OSPF router in the traffic-engineered portion of your network.

SUMMARY STEPS

1. **configure**
2. **router ospf** *process-name*
3. **mpls traffic-eng router-id** *interface*
4. **segment-routing mpls**
5. **area** *area*
6. **mpls traffic-eng**
7. **segment-routing mpls**
8. **exit**
9. **mpls traffic-eng**
10. Use the **commit** or **end** command.

DETAILED STEPS

	Command or Action	Purpose
Step 1	**configure** **Example:** `RP/0/RP0/CPU0:router# configure`	Enters mode.
Step 2	**router ospf** *process-name* **Example:** `RP/0/RP0/CPU0:router(config)# router ospf 1`	Enables OSPF routing for the specified routing process and places the router in router configuration mode.
Step 3	**mpls traffic-eng router-id** *interface* **Example:** `RP/0/RP0/CPU0:router(config-ospf)# mpls traffic-eng router-id Loopback0`	Sets the traffic engineering loopback interface.
Step 4	**segment-routing mpls** **Example:** `RP/0/RP0/CPU0:router(config-ospf)# segment-routing mpls`	Enables segment routing using the MPLS data plane on the routing process and all areas and interfaces in the routing process. Enables segment routing fowarding on all interfaces in the routing process and installs the SIDs received by OSPF in the forwarding table.
Step 5	**area** *area* **Example:** `RP/0/RP0/CPU0:router(config-ospf)# area 0`	Enters area configuration mode.

	Command or Action	Purpose
Step 6	**mpls traffic-eng** **Example:** `RP/0/RP0/CPU0:router(config-ospf-ar)# mpls traffic-eng`	Enables IGP traffic engineering funtionality.
Step 7	**segment-routing mpls** **Example:** `RP/0/RP0/CPU0:router(config-ospf-ar)# segment-routing mpls`	(Optional) Enables segment routing using the MPLS data plane on the area and all interfaces in the area. Enables segment routing fowarding on all interfaces in the area and installs the SIDs received by OSPF in the forwarding table.

Step 8	**exit**
	Example:
	`RP/0/RP0/CPU0:router(config-ospf-ar)# exit`
	`RP/0/RP0/CPU0:router(config-ospf)# exit`	

Step 9	**mpls traffic-eng** **Example:** `RP/0/RP0/CPU0:router(config)# mpls traffic-eng`	Enables traffic engineering funtionality on the node. The node advertises the traffic engineering link attributes in IGP which populates the traffic engineering database (TED) on the head-end. The SR-TE head-end requires the TED to calculate and validate the path of the SR-TE policy.
Step 10	Use the **commit** or **end** command.	**commit** —Saves the configuration changes and remains within the configuration session. **end** —Prompts user to take one of these actions: • **Yes** — Saves configuration changes and exits the configuration session. • **No** —Exits the configuration session without committing the configuration changes. • **Cancel** —Remains in the configuration session, without committing the configuration changes.

What to do next

Configure the prefix SID.

Configuring a Prefix-SID on the OSPF-Enabled Loopback Interface

A prefix segment identifier (SID) is associated with an IP prefix. The prefix SID is manually configured from the segment routing global block (SRGB) range of labels. A prefix SID is configured under the loopback interface with the loopback address of the node as the prefix. The prefix segment steers the traffic along the shortest path to its destination.

A prefix SID can be a node SID or an Anycast SID. A node SID is a type of prefix SID that identifies a specific node. An Anycast SID is a type of prefix SID that identifies a set of nodes, and is configured with n-flag clear. The set of nodes (Anycast group) is configured to advertise a shared prefix address and prefix SID. Anycast routing enables the steering of traffic toward multiple advertising nodes. Packets addressed to an Anycast address are forwarded to the topologically nearest nodes.

The prefix SID is globally unique within the segment routing domain.

This task describes how to configure prefix segment identifier (SID) index or absolute value on the OSPF-enabled Loopback interface.

Before you begin

Ensure that segment routing is enabled on an instance, area, or interface.

SUMMARY STEPS

1. **configure**
2. **router ospf** *process-name*
3. **area** *value*
4. **interface Loopback** *interface-instance*
5. **prefix-sid**{**index** *SID-index* | **absolute** *SID-value* } [**n-flag-clear**] [**explicit-null**]
6. Use the **commit** or **end** command.

DETAILED STEPS

	Command or Action	Purpose
Step 1	**configure** **Example:** `RP/0/RP0/CPU0:router# configure`	Enters mode.
Step 2	**router ospf** *process-name* **Example:** `RP/0/RP0/CPU0:router(config)# router ospf 1`	Enables OSPF routing for the specified routing process, and places the router in router configuration mode.
Step 3	**area** *value* **Example:** `RP/0/RP0/CPU0:router(config-ospf)# area 0`	Enters area configuration mode.
Step 4	**interface Loopback** *interface-instance* **Example:** `RP/0/RP0/CPU0:router(config-ospf-ar)# interface Loopback0 passive`	Specifies the loopback interface and instance.
Step 5	**prefix-sid**{**index** *SID-index* | **absolute** *SID-value* } [**n-flag-clear**] [**explicit-null**]	Configures the prefix-SID index or absolute value for the interface.

	Command or Action	Purpose

	Example:
	`RP/0/RP0/CPU0:router(config-ospf-ar)# prefix-sid index 1001`
	`RP/0/RP0/CPU0:router(config-ospf-ar)# prefix-sid absolute 17001`
		Specify **index** *SID-index* for each node to create a prefix SID based on the lower boundary of the SRGB + the index.
		Specify **absolute** *SID-value* for each node to create a specific prefix SID within the SRGB.
		By default, the n-flag is set on the prefix-SID, indicating that it is a node SID. For specific prefix-SID (for example, Anycast prefix-SID), enter the `n-flag-clear` keyword. OSPF does not set the `N` flag in the prefix-SID sub Type Length Value (TLV).
		To disable penultimate-hop-popping (PHP) and add an explicit-Null label, enter the `explicit-null` keyword. OSPF sets the `E` flag in the prefix-SID sub TLV.

Step 6	Use the **commit** or **end** command.	**commit** —Saves the configuration changes and remains within the configuration session. **end** —Prompts user to take one of these actions: • **Yes** — Saves configuration changes and exits the configuration session. • **No** —Exits the configuration session without committing the configuration changes. • **Cancel** —Remains in the configuration session, without committing the configuration changes.

Verify the prefix-SID configuration:

```
RP/0/RP0/CPU0:router# show ospf database opaque-area 7.0.0.1 self-originate
 OSPF Router with ID (10.0.0.1) (Process ID 1)
                Type-10 Opaque Link Area Link States (Area 0)
<...>
    Extended Prefix TLV: Length: 20
      Route-type: 1
      AF         : 0
      Flags      : 0x40
      Prefix     : 10.0.0.1/32

      SID sub-TLV: Length: 8
        Flags      : 0x0
        MTID       : 0
        Algo       : 0
        SID Index  : 1001
```

What to do next

Configure SR-TE Policies

CHAPTER 6

Configure Segment Routing for BGP

Border Gateway Protocol (BGP) is an Exterior Gateway Protocol (EGP) that allows you to create loop-free inter-domain routing between autonomous systems. An autonomous system is a set of routers under a single technical administration. Routers in an autonomous system can use multiple Interior Gateway Protocols (IGPs) to exchange routing information inside the autonomous system and an EGP to route packets outside the autonomous system.

This module provides the configuration information used to enable Segment Routing for BGP.

Note For additional information on implementing BGP on your router, see the *BGP Configuration Guide for Cisco NCS 5500 Series Routers*.

Segment Routing for BGP

In a traditional BGP-based data center (DC) fabric, packets are forwarded hop-by-hop to each node in the autonomous system. Traffic is directed only along the external BGP (eBGP) multipath ECMP. No traffic engineering is possible.

In an MPLS-based DC fabric, the eBGP sessions between the nodes exchange BGP labeled unicast (BGP-LU) network layer reachability information (NLRI). An MPLS-based DC fabric allows any leaf (top-of-rack or border router) in the fabric to communicate with any other leaf using a single label, which results in higher packet forwarding performance and lower encapsulation overhead than traditional BGP-based DC fabric. However, since each label value might be different for each hop, an MPLS-based DC fabric is more difficult to troubleshoot and more complex to configure.

BGP has been extended to carry segment routing prefix-SID index. BGP-LU helps each node learn BGP prefix SIDs of other leaf nodes and can use ECMP between source and destination. Segment routing for BGP simplifies the configuration, operation, and troubleshooting of the fabric. With segment routing for BGP, you can enable traffic steering capabilities in the data center using a BGP prefix SID.

Configure BGP Prefix Segment Identifiers

Segments associated with a BGP prefix are known as BGP prefix SIDs. The BGP prefix SID is global within a segment routing or BGP domain. It identifies an instruction to forward the packet over the ECMP-aware best-path computed by BGP to the related prefix. The BGP prefix SID is manually configured from the segment routing global block (SRGB) range of labels.

Each BGP speaker must be configured with an SRGB using the **segment-routing global-block** command. See the About the Segment Routing Global Block section for information about the SRGB.

Note Because the values assigned from the range have domain-wide significance, we recommend that all routers within the domain be configured with the same range of values.

To assign a BGP prefix SID, first create a routing policy using the **set label-index** *index* attribute, then associate the index to the node.

Note A routing policy with the **set label-index** attribute can be attached to a network configuration or redistribute configuration. Other routing policy language (RPL) configurations are possible. For more information on routing policies, refer to the "Implementing Routing Policy" chapter in the *Routing Configuration Guide for Cisco NCS 5500 Series Routers*.

Example

The following example shows how to configure the SRGB, create a BGP route policy using a $SID parameter and **set label-index** attribute, and then associate the prefix-SID index to the node.

```
RP/0/RSP0/CPU0:router(config)# segment-routing global-block 16000 23999

RP/0/RSP0/CPU0:router(config)# route-policy SID($SID)
RP/0/RSP0/CPU0:router(config-rpl)# set label-index $SID
RP/0/RSP0/CPU0:router(config-rpl)# end policy

RP/0/RSP0/CPU0:router(config)# router bgp 1
RP/0/RSP0/CPU0:router(config-bgp)# bgp router-id 1.1.1.1
RP/0/RSP0/CPU0:router(config-bgp)# address-family ipv4 unicast
RP/0/RSP0/CPU0:router(config-bgp-af)# network 1.1.1.3/32 route-policy SID(3)
RP/0/RSP0/CPU0:router(config-bgp-af)# allocate-label all
RP/0/RSP0/CPU0:router(config-bgp-af)# commit
RP/0/RSP0/CPU0:router(config-bgp-af)# end


RP/0/RSP0/CPU0:router# show bgp 1.1.1.3/32
BGP routing table entry for 1.1.1.3/32
Versions:
  Process           bRIB/RIB  SendTblVer
  Speaker                 74          74
    Local Label: 16003
Last Modified: Sep 29 19:52:18.155 for 00:07:22
Paths: (1 available, best #1)
  Advertised to update-groups (with more than one peer):
    0.2
```

```
Path #1: Received by speaker 0
Advertised to update-groups (with more than one peer):
  0.2
3
  99.3.21.3 from 99.3.21.3 (1.1.1.3)
    Received Label 3
    Origin IGP, metric 0, localpref 100, valid, external, best, group-best
    Received Path ID 0, Local Path ID 1, version 74
    Origin-AS validity: not-found
    Label Index: 3
```

Segment Routing Egress Peer Engineering

Segment routing egress peer engineering (EPE) uses a controller to instruct an ingress provider edge, or a content source (node) within the segment routing domain, to use a specific egress provider edge (node) and a specific external interface to reach a destination. BGP peer SIDs are used to express source-routed inter-domain paths.

Below are the BGP-EPE peering SID types:

- PeerNode SID—To an eBGP peer. Pops the label and forwards the traffic on any interface to the peer.
- PeerAdjacency SID—To an eBGP peer via interface. Pops the label and forwards the traffic on the related interface.

The controller learns the BGP peer SIDs and the external topology of the egress border router through BGP-LS EPE routes. The controller can program an ingress node to steer traffic to a destination through the egress node and peer node using BGP labeled unicast (BGP-LU).

EPE functionality is only required at the EPE egress border router and the EPE controller.

Configure Segment Routing Egress Peer Engineering

This task explains how to configure segment routing EPE on the EPE egress node.

SUMMARY STEPS

1. **router bgp** *as-number*
2. **neighbor** *ip-address*
3. **remote-as** *as-number*
4. **egress-engineering**

DETAILED STEPS

	Command or Action	Purpose
Step 1	**router bgp** *as-number* **Example:** `RP/0/RSP0/CPU0:router(config)# router bgp 1`	Specifies the BGP AS number and enters the BGP configuration mode, allowing you to configure the BGP routing process.

	Command or Action	Purpose
Step 2	**neighbor** *ip-address* **Example:** `RP/0/RSP0/CPU0:router(config-bgp)# neighbor 192.168.1.3`	Places the router in neighbor configuration mode for BGP routing and configures the neighbor IP address as a BGP peer.
Step 3	**remote-as** *as-number* **Example:** `RP/0/RSP0/CPU0:router(config-bgp-nbr)# remote-as 3`	Creates a neighbor and assigns a remote autonomous system number to it.
Step 4	**egress-engineering** **Example:** `RP/0/RSP0/CPU0:router(config-bgp-nbr)# egress-engineering`	Configures the egress node with EPE for the eBGP peer.

Configure BGP Link-State

BGP Link-State (LS) is an Address Family Identifier (AFI) and Sub-address Family Identifier (SAFI) defined to carry interior gateway protocol (IGP) link-state database through BGP. BGP LS delivers network topology information to topology servers and Application Layer Traffic Optimization (ALTO) servers. BGP LS allows policy-based control to aggregation, information-hiding, and abstraction. BGP LS supports IS-IS and OSPFv2.

Note IGPs do not use BGP LS data from remote peers. BGP does not download the received BGP LS data to any other component on the router.

For segment routing, the following attributes have been added to BGP LS:

- Node—Segment routing capability (including SRGB range) and algorithm
- Link—Adjacency SID and LAN adjacency SID
- Prefix—Prefix SID and segment routing mapping server (SRMS) prefix range

The following example shows how to exchange link-state information with a BGP neighbor:

```
RP/0/RSP0/CPU0:router# configure
RP/0/RSP0/CPU0:router(config)# router bgp 1
RP/0/RSP0/CPU0:router(config-bgp)# neighbor 10.0.0.2
RP/0/RSP0/CPU0:router(config-bgp-nbr)# remote-as 1
RP/0/RSP0/CPU0:router(config-bgp-nbr)# address-family link-state link-state
RP/0/RSP0/CPU0:router(config-bgp-nbr-af)# exit
```

IGP Link-State Database Distribution

A given BGP node may have connections to multiple, independent routing domains. IGP link-state database distribution into BGP-LS is supported for both OSPF and IS-IS protocols in order to distribute this information on to controllers or applications that desire to build paths spanning or including these multiple domains.

To distribute IS-IS link-state data using BGP LS, use the **distribute link-state** command in router configuration mode.

```
RP/0/RSP0/CPU0:router# configure
RP/0/RSP0/CPU0:router(config)# router isis isp
RP/0/RSP0/CPU0:router(config-isis)# distribute link-state instance-id 32 level 2 throttle 
5
```

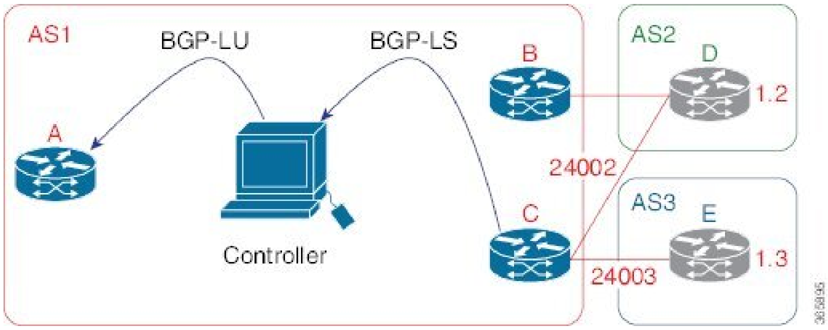
To distribute OSPFv2 link-state data using BGP LS, use the **distribute link-state** command in router configuration mode.

```
RP/0/RSP0/CPU0:router# configure
RP/0/RSP0/CPU0:router(config)# router ospf 100
RP/0/RSP0/CPU0:router(config-ospf)# distribute link-state instance-id 32 throttle 10
```

Example: Configuring SR-EPE and BGP-LS

In the following figure, segment routing is enabled on autonomous system AS1 with ingress node A and egress nodes B and C. In this example, we configure EPE on egress node C.

Figure 2: Topology



Step 1 Configure node C with EPE for eBGP peers D and E.

Example:

```
RP/0/RSP0/CPU0:router_C(config)# router bgp 1
RP/0/RSP0/CPU0:router_C(config-bgp)# neighbor 192.168.1.3
RP/0/RSP0/CPU0:router_C(config-bgp-nbr)# remote-as 3
RP/0/RSP0/CPU0:router_C(config-bgp-nbr)# description to E
RP/0/RSP0/CPU0:router_C(config-bgp-nbr)# egress-engineering
RP/0/RSP0/CPU0:router_C(config-bgp-nbr)# address-family ipv4 unicast
RP/0/RSP0/CPU0:router_C(config-bgp-nbr-af)# route-policy bgp_in in
RP/0/RSP0/CPU0:router_C(config-bgp-nbr-af)# route-policy bgp_out out
RP/0/RSP0/CPU0:router_C(config-bgp-nbr-af)# exit
```

```
RP/0/RSP0/CPU0:router_C(config-bgp-nbr)# exit
RP/0/RSP0/CPU0:router_C(config-bgp)# neighbor 192.168.1.2
RP/0/RSP0/CPU0:router_C(config-bgp-nbr)# remote-as 2
RP/0/RSP0/CPU0:router_C(config-bgp-nbr)# description to D
RP/0/RSP0/CPU0:router_C(config-bgp-nbr)# egress-engineering
RP/0/RSP0/CPU0:router_C(config-bgp-nbr)# address-family ipv4 unicast
RP/0/RSP0/CPU0:router_C(config-bgp-nbr-af)# route-policy bgp_in in
RP/0/RSP0/CPU0:router_C(config-bgp-nbr-af)# route-policy bgp_out out
RP/0/RSP0/CPU0:router_C(config-bgp-nbr-af)# exit
RP/0/RSP0/CPU0:router_C(config-bgp-nbr)# exit
```

Step 2 Configure node C to advertise peer node SIDs to the controller using BGP-LS.

Example:

```
RP/0/RSP0/CPU0:router_C(config-bgp)# neighbor 172.29.50.71
RP/0/RSP0/CPU0:router_C(config-bgp-nbr)# remote-as 1
RP/0/RSP0/CPU0:router_C(config-bgp-nbr)# description to EPE_controller
RP/0/RSP0/CPU0:router_C(config-bgp-nbr)# address-family link-state link-state
RP/0/RSP0/CPU0:router_C(config-bgp-nbr)# exit
RP/0/RSP0/CPU0:router_C(config-bgp)# exit
```

Step 3 Commit the configuration.

Example:

```
RP/0/RSP0/CPU0:router_C(config)# commit
```

Step 4 Verify the configuration.

Example:

```
RP/0/RSP0/CPU0:router_C# show bgp egress-engineering

 Egress Engineering Peer Set: 192.168.1.2/32 (10b87210)
     Nexthop: 192.168.1.2
     Version: 2, rn_version: 2
       Flags: 0x00000002
   Local ASN: 1
  Remote ASN: 2
   Local RID: 1.1.1.3
  Remote RID: 1.1.1.4
   First Hop: 192.168.1.2
        NHID: 3
       Label: 24002, Refcount: 3
     rpc_set: 10b9d408

 Egress Engineering Peer Set: 192.168.1.3/32 (10be61d4)
     Nexthop: 192.168.1.3
     Version: 3, rn_version: 3
       Flags: 0x00000002
   Local ASN: 1
  Remote ASN: 3
   Local RID: 1.1.1.3
  Remote RID: 1.1.1.5
   First Hop: 192.168.1.3
        NHID: 4
       Label: 24003, Refcount: 3
     rpc_set: 10be6250
```

The output shows that node C has allocated peer SIDs for each eBGP peer.

Example:

```
RP/0/RSP0/CPU0:router_C# show mpls forwarding labels 24002 24003
Local  Outgoing    Prefix             Outgoing     Next Hop        Bytes
Label  Label       or ID              Interface                    Switched
------ ----------- ------------------ ------------ --------------- ------------
24002  Unlabelled  No ID              Te0/3/0/0    192.168.1.2     0
24003  Unlabelled  No ID              Te0/1/0/0    192.168.1.3     0
```

The output shows that node C installed peer node SIDs in the Forwarding Information Base (FIB).

CHAPTER 7

Configure SR-TE Policies

This module provides information about segment routing for traffic engineering (SR-TE) policies, how to configure SR-TE policies, and how to steer traffic into an SR-TE policy.

Note Configuring SR-TE policies with 3 or more labels and an L2 Transport Interface on the same network processing unit (NPU) can cause traffic loss.

About SR-TE Policies

Segment routing for traffic engineering (SR-TE) uses a “policy” to steer traffic through the network. An SR-TE policy path is expressed as a list of segments that specifies the path, called a segment ID (SID) list. Each segment is an end-to-end path from the source to the destination, and instructs the routers in the network to follow the specified path instead of the shortest path calculated by the IGP. If a packet is steered into an SR-TE policy, the SID list is pushed on the packet by the head-end. The rest of the network executes the instructions embedded in the SID list.

There are two types of SR-TE policies: dynamic and explicit.

Local Dynamic SR-TE Policy

When you configure local dynamic SR-TE, the head-end locally calculates the path to the destination address. Dynamic path calculation results in a list of interface IP addresses that traffic engineering (TE) maps to adj-SID labels. Routes are learned by way of forwarding adjacencies over the TE tunnel.

Explicit SR-TE Policy

An explicit path is a list of IP addresses or labels, each representing a node or link in the explicit path. This feature is enabled through the **explicit-path** command that allows you to create an explicit path and enter a configuration submode for specifying the path.

How to Configure SR-TE Policies

This section contains the following procedures:

- Configure Local Dynamic SR-TE Policy, on page 34
- Configure Explicit SR-TE Policy, on page 35

Configure Local Dynamic SR-TE Policy

This task explains how to configure a local dynamic SR-TE policy.

SUMMARY STEPS

1. **configure**
2. **interface tunnel-te** *tunnel-id*
3. **ipv4 unnumbered** *type interface-path-id*
4. **destination** *ip-address*
5. **path-option** *preference-priority* **dynamic segment-routing**
6. **path-protection**
7. **commit**

DETAILED STEPS

	Command or Action	Purpose
Step 1	**configure** **Example:** `RP/0/RP0/CPU0:router# configure`	Enters mode.
Step 2	**interface tunnel-te** *tunnel-id* **Example:** `RP/0/RP0/CPU0:router(config)# interface tunnel-te22`	Configures the tunnel interface.
Step 3	**ipv4 unnumbered** *type interface-path-id* **Example:** `RP/0/RP0/CPU0:router(config-if)# ipv4 unnumbered loopback0`	Assigns a source address so that forwarding can be performed on the new tunnel. Loopback is commonly used as the interface type.
Step 4	**destination** *ip-address* **Example:** `RP/0/RP0/CPU0:router(config-if)# destination`	Assigns a destination address on the new tunnel.

	Command or Action	Purpose
	`192.168.0.2`	
Step 5	**path-option** *preference-priority* **dynamic segment-routing** **Example:** `RP/0/RP0/CPU0:router(config-if)# path-option 1 dynamic segment-routing`	Sets the path option to dynamic and assigns the path ID.
Step 6	**path-protection** **Example:** `RP/0/RP0/CPU0:router(config-if)# path-protection`	Enables path protection on the tunnel-te interface.
Step 7	**commit**	

This completes the configuration of the dynamic SR-TE policy.

Configure Explicit SR-TE Policy

This task explains how to configure an explicit SR-TE policy.

SUMMARY STEPS

1. **configure**
2. **explicit-path name** *path-name*
3. **index** *index* {**next-address** *ip-address* | **next-label** *label*}
4. **exit**
5. **interface tunnel-te** *tunnel-id*
6. **ipv4 unnumbered** *type interface-path-id*
7. **destination** *ip-address* [**verbatim**]
8. **path-option** *preference-priority* **explicit name** *path-name* **segment-routing**
9. **commit**

DETAILED STEPS

	Command or Action	Purpose
Step 1	**configure** **Example:** `RP/0/RP0/CPU0:router# configure`	Enters mode.
Step 2	**explicit-path name** *path-name* **Example:**	Enters a name for the explicit path and enters the explicit path configuration mode.

	Command or Action	Purpose
	`RP/0/RP0/CPU0:router(config)# explicit-path name r1r6_exp`	

Step 3	**index** *index* {**next-address** *ip-address* | **next-label** *label*}
	Example:
	`RP/0/RP0/CPU0:router(config-expl-path)# index 1 next-label 16001`
	`RP/0/RP0/CPU0:router(config-expl-path)# index 2 next-label 16006`
		Specifies a label or an address in an explicit path of a tunnel.
		Note
		- You can include multiple addresses, labels, or both. However, you cannot configure addresses after you have configured labels. Once you start configuring labels, you need to continue with labels.
		- Each entry must have a unique index.
		- If the first hop is specified as next-label, that label must be an Adj-SID of the head-end or a prefix-SID label value known by the head-end.

Step 4	**exit**	
Step 5	**interface tunnel-te** *tunnel-id* **Example:** `RP/0/RP0/CPU0:router(config)# interface tunnel-te22`	Configures the tunnel interface.
Step 6	**ipv4 unnumbered** *type interface-path-id* **Example:** `RP/0/RP0/CPU0:router(config-if)# ipv4 unnumbered loopback0`	Assigns a source address so that forwarding can be performed on the new tunnel. Loopback is commonly used as the interface type.
Step 7	**destination** *ip-address* [**verbatim**] **Example:** `RP/0/RP0/CPU0:router(config-if)# destination 192.168.0.2`	Assigns a destination address on the new tunnel. Typically, the tunnel destination must have a match in the routing information base (RIB). For inter-area or inter-domain policies to destinations that are otherwise not reachable, use the **verbatim** option to disable the RIB verification on a tunnel destination.
Step 8	**path-option** *preference-priority* **explicit name** *path-name* **segment-routing** **Example:** `RP/0/RP0/CPU0:router(config-if)# path-option 1 explicit name r1r6_exp segment-routing`	Specifies the explicit path name and assigns the path ID.
Step 9	**commit**	

This completes the configuration of the explicit SR-TE policy.

Steering Traffic into an SR-TE Policy

This section describes the following traffic steering methods:

Static Routes

Static routes can use the segment routing tunnel as a next-hop interface. Both IPv4 and IPv6 prefixes can be routed through the tunnel.

A static route to a destination with a prefix-SID removes the IGP-installed SR-forwarding entry of that prefix.

Autoroute Announce

The SR-TE policy can be advertised into an IGP as a next hop by configuring the autoroute announce statement on the source router. The IGP then installs routes in the Routing Information Base (RIB) for shortest paths that involve the tunnel destination. Autoroute announcement of IPv4 prefixes can be carried through either OSPF or IS-IS. Autoroute announcement of IPv6 prefixes can be carried only through IS-IS.

Autoroute Destination

Autoroute destination allows you to automatically route traffic through a segment routing tunnel instead of manually configuring static routes. Multiple autoroute destination addresses can be added in the routing information base (RIB) per tunnel.

Static routes are always added with zero cost metric, which can result in traffic that is mapped on multiple tunnels to always load-balance due to ECMP. This load-balancing may be undesirable when some of those tunnels have sub-optimal paths. With autoroute destination, only the tunnel whose IGP cost to its endpoint is lowest will be considered for carrying traffic.

- **Interaction Between Static Routes and Autoroute Destination**

  If there is a manually configured static route to the same destination as a tunnel with autoroute destination enabled, traffic for that destination is load-shared between the static route and the tunnel with autoroute destination enabled.

- **Interaction Between Autoroute Announce and Autoroute Destination**

  For intra-area tunnels, if a tunnel is configured with both autoroute announce and autoroute destination, the tunnel is announced to the RIB by both the IGP and the static process. RIBs prefer static routes, not IGP routes, so the autoroute destination features takes precedence over autoroute announce.

Configure Static Routes

This task explains how to configure a static route.

SUMMARY STEPS

1. **configure**
2. **interface tunnel-te** *tunnel-id*
3. **ipv4 unnumbered** *type interface-path-id*

4. **destination** *ip-address*
5. **path-option** *preference-priority* **dynamic segment-routing**
6. **exit**
7. **router static**
8. **address-family ipv4 unicast**
9. *prefix mask interface-type interface-instance*
10. **commit**

DETAILED STEPS

	Command or Action	Purpose
Step 1	**configure** **Example:** `RP/0/RP0/CPU0:router# configure`	Enters mode.
Step 2	**interface tunnel-te** *tunnel-id* **Example:** `RP/0/RP0/CPU0:router(config)# interface tunnel-te22`	Configures the tunnel interface.
Step 3	**ipv4 unnumbered** *type interface-path-id* **Example:** `RP/0/RP0/CPU0:router(config-if)# ipv4 unnumbered loopback0`	Assigns a source address so that forwarding can be performed on the new tunnel. Loopback is commonly used as the interface type.
Step 4	**destination** *ip-address* **Example:** `RP/0/RP0/CPU0:router(config-if)# destination 192.168.0.2`	Assigns a destination address on the new tunnel.
Step 5	**path-option** *preference-priority* **dynamic segment-routing** **Example:** `RP/0/RP0/CPU0:router(config-if)# path-option 1 dynamic segment-routing`	Sets the path option to dynamic and assigns the path ID.
Step 6	**exit**	
Step 7	**router static** **Example:**	Configures the static route and enters static configuration mode.

	Command or Action	Purpose
	`RP/0/RP0/CPU0:router(config)# router static`	
Step 8	**address-family ipv4 unicast** **Example:** `RP/0/RP0/CPU0:router(config-static)# address-family ipv4 unicast`	Enters address family mode.
Step 9	*prefix mask interface-type interface-instance* **Example:** `RP/0/RP0/CPU0:router(config-static-af)# 192.168.0.2/32 tunnel-te22`	Specifies the destination prefix is directly reachable through the tunnel interface.
Step 10	**commit**	

This completes the configuration of the static route.

Configure Autoroute Announce

This task explains how to configure autoroute announce to steer traffic through the SR-TE policy.

SUMMARY STEPS

1. **configure**
2. **interface tunnel-te** *tunnel-id*
3. **ipv4 unnumbered** *type interface-path-id*
4. **autoroute announce**
5. **destination** *ip-address*
6. **path-option** *preference-priority* **dynamic segment-routing**
7. **path-protection**
8. **commit**

DETAILED STEPS

	Command or Action	Purpose
Step 1	**configure** **Example:** `RP/0/RP0/CPU0:router# configure`	Enters mode.
Step 2	**interface tunnel-te** *tunnel-id* **Example:**	Configures the tunnel interface.

	Command or Action	Purpose
	`RP/0/RP0/CPU0:router(config)# interface tunnel-te22`	
Step 3	**ipv4 unnumbered** *type interface-path-id* **Example:** `RP/0/RP0/CPU0:router(config-if)# ipv4 unnumbered loopback0`	Assigns a source address so that forwarding can be performed on the new tunnel. Loopback is commonly used as the interface type.
Step 4	**autoroute announce** **Example:** `RP/0/RP0/CPU0:router(config-if)# autoroute announce`	Enables messages that notify the neighbor nodes about the routes that are forwarding.
Step 5	**destination** *ip-address* **Example:** `RP/0/RP0/CPU0:router(config-if)# destination 192.168.0.2`	Assigns a destination address on the new tunnel.
Step 6	**path-option** *preference-priority* **dynamic segment-routing** **Example:** `RP/0/RP0/CPU0:router(config-if)# path-option 1 dynamic segment-routing`	Sets the path option to dynamic and assigns the path ID.
Step 7	**path-protection** **Example:** `RP/0/RP0/CPU0:router(config-if)# path-protection`	Enables path protection on the tunnel-te interface.
Step 8	**commit**	

Configure Autoroute Destination

This task explains how to configure autoroute destination to steer traffic through the SR-TE policy.

SUMMARY STEPS

1. **configure**
2. **interface tunnel-te** *tunnel-id*
3. **ipv4 unnumbered** *type interface-path-id*

4. **autoroute destination** *destination-ip-address*
5. **destination** *ip-address*
6. **path-option** *preference-priority* **dynamic segment-routing**
7. **commit**

DETAILED STEPS

	Command or Action	Purpose
Step 1	**configure** **Example:** `RP/0/RP0/CPU0:router# configure`	Enters mode.
Step 2	**interface tunnel-te** *tunnel-id* **Example:** `RP/0/RP0/CPU0:router(config)# interface tunnel-te22`	Configures the tunnel interface.
Step 3	**ipv4 unnumbered** *type interface-path-id* **Example:** `RP/0/RP0/CPU0:router(config-if)# ipv4 unnumbered loopback0`	Assigns a source address so that forwarding can be performed on the new tunnel. Loopback is commonly used as the interface type.

Step 4	**autoroute destination** *destination-ip-address*
	Example:
	`RP/0/RP0/CPU0:router(config-if)# autoroute destination 192.168.0.1`
	`RP/0/RP0/CPU0:router(config-if)# autoroute destination 192.168.0.2  (the default route)`
	`RP/0/RP0/CPU0:router(config-if)# autoroute destination 192.168.0.3`
	`RP/0/RP0/CPU0:router(config-if)# autoroute destination 192.168.0.4`
		(Optional) Adds a route (destination-ip-address) in the RIB with the tunnel as outgoing interface to the tunnel destination.

Step 5	**destination** *ip-address* **Example:** `RP/0/RP0/CPU0:router(config-if)# destination 192.168.0.2`	Assigns a destination address on the new tunnel.
Step 6	**path-option** *preference-priority* **dynamic segment-routing** **Example:** `RP/0/RP0/CPU0:router(config-if)# path-option 1`	Sets the path option to dynamic and assigns the path ID.

	Command or Action	Purpose
	`dynamic segment-routing`	
Step 7	**commit**	

BGP SR-TE

SR-TE can be used by data center (DC) operators to provide different levels of Service Level Assurance (SLA). Setting up SR-TE paths using BGP (BGP SR-TE) simplifies DC network operation without introducing a new protocol for this purpose.

Explicit BGP SR-TE

Explicit BGP SR-TE uses an SR-TE policy (identified by a unique color ID) that contains a list of explicit paths with SIDs that correspond to each explicit path. A BGP speaker signals an explicit SR-TE policy to a remote peer, which triggers the setup of a TE tunnel with specific characteristics and explicit paths. On the receiver side, a TE tunnel that corresponds to the explicit path is setup by BGP. The packets for the destination mentioned in the BGP update follow the explicit path described by the policy. Each policy can include multiple explicit paths, and TE will create a tunnel for each path.

Note For more information on routing policies and routing policy language (RPL), refer to the "Implementing Routing Policy" chapter in the *Routing Configuration Guide for Cisco NCS 5500 Series Routers*.

IPv4 and IPv6 SR policies can be advertised over BGPv4 or BGPv6 sessions between the SR-TE controller and the SR-TE headend. The Cisco IOS-XR implementation supports the following combinations:

- IPv4 SR policy advertised over BGPv4 session
- IPv6 SR policy advertised over BGPv4 session
- IPv6 SR policy advertised over BGPv6 session

Configure Explicit BGP SR-TE

Perform this task to configure explicit BGP SR-TE:

SUMMARY STEPS

1. **configure**
2. **router bgp** *as-number*
3. **bgp router-id** *ip-address*
4. **address-family** {**ipv4** | **ipv6**} **sr-policy**
5. **exit**
6. **neighbor** *ip-address*
7. **remote-as** *as-number*
8. **address-family** {**ipv4** | **ipv6**} **sr-policy**

9. **route-policy** *route-policy-name* {**in** | **out**}

DETAILED STEPS

	Command or Action	Purpose
Step 1	**configure**	
Step 2	**router bgp** *as-number* **Example:** `RP/0/RSP0/CPU0:router(config)# router bgp 65000`	Specifies the BGP AS number and enters the BGP configuration mode, allowing you to configure the BGP routing process.
Step 3	**bgp router-id** *ip-address* **Example:** `RP/0/RSP0/CPU0:router(config-bgp)# bgp router-id 1.1.1.1`	Configures the local router with a specified router ID.
Step 4	**address-family** {**ipv4** | **ipv6**} **sr-policy** **Example:** `RP/0/RSP0/CPU0:router(config-bgp)# address-family ipv4 sr-policy`	Specifies either the IPv4 or IPv6 address family and enters address family configuration submode.
Step 5	**exit**	
Step 6	**neighbor** *ip-address* **Example:** `RP/0/RSP0/CPU0:router(config-bgp)# neighbor 10.10.0.1`	Places the router in neighbor configuration mode for BGP routing and configures the neighbor IP address as a BGP peer.
Step 7	**remote-as** *as-number* **Example:** `RP/0/RSP0/CPU0:router(config-bgp-nbr)# remote-as 1`	Creates a neighbor and assigns a remote autonomous system number to it.
Step 8	**address-family** {**ipv4** | **ipv6**} **sr-policy** **Example:** `RP/0/RSP0/CPU0:router(config-bgp-nbr)# address-family ipv4 sr-policy`	Specifies either the IPv4 or IPv6 address family and enters address family configuration submode.
Step 9	**route-policy** *route-policy-name* {**in** | **out**} **Example:**	Applies the specified policy to IPv4 or IPv6 unicast routes.

	Command or Action	Purpose
	`RP/0/RSP0/CPU0:router(config-bgp-nbr-af)#` **`route-policy pass out`**	

Example: BGP SR-TE with BGPv4 Neighbor to BGP SR-TE Controller

The following configuration shows the an SR-TE head-end with a BGPv4 session towards a BGP SR-TE controller. This BGP session is used to signal both IPv4 and IPv6 SR policies.

```
router bgp 65000
bgp router-id 1.1.1.1
 !
 address-family ipv4 sr-policy
 !
 address-family ipv6 sr-policy
 !
 neighbor 10.1.3.1
  remote-as 10
  description *** eBGP session to BGP SRTE controller ***
  address-family ipv4 sr-policy
   route-policy pass in
   route-policy pass out
  !
  address-family ipv6 sr-policy
   route-policy pass in
   route-policy pass out
  !
 !
!
```

Example: BGP SR-TE with BGPv6 Neighbor to BGP SR-TE Controller

The following configuration shows an SR-TE head-end with a BGPv6 session towards a BGP SR-TE controller. This BGP session is used to signal IPv6 SR policies.

```
router bgp 65000
bgp router-id 1.1.1.1
 address-family ipv6 sr-policy
 !
 neighbor 3001::10:1:3:1
  remote-as 10
  description *** eBGP session to BGP SRTE controller ***
  address-family ipv6 sr-policy
   route-policy pass in
   route-policy pass out
  !
 !
!
```

Using Binding Segments

The binding segment is a local segment identifying an SR-TE policy. Each SR-TE policy is associated with a binding segment ID (BSID). The BSID is a local label that is automatically allocated for each SR-TE policy when the SR-TE policy is instantiated.

BSID can be used to steer traffic into the SR-TE policy and across domain borders, creating seamless end-to-end inter-domain SR-TE policies. Each domain controls its local SR-TE policies; local SR-TE policies can be validated and rerouted if needed, independent from the remote domain's head-end. Using binding segments isolates the head-end from topology changes in the remote domain.

Packets received with a BSID as top label are steered into the SR-TE policy associated with the BSID. When the BSID label is popped, the SR-TE policy's SID list is pushed.

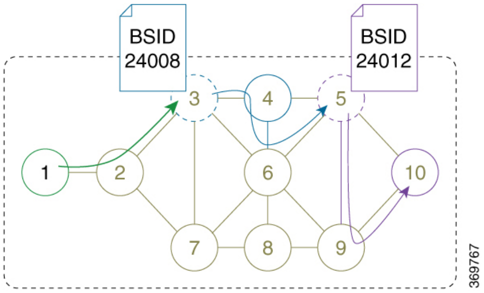
BSID can be used in the following cases:

- Multi-Domain (inter-domain, inter-autonomous system)—BSIDs can be used to steer traffic across domain borders, creating seamless end-to-end inter-domain SR-TE policies.
- Large-Scale within a single domain—The head-end can use hierarchical SR-TE policies by nesting the end-to-end (edge-to-edge) SR-TE policy within another layer of SR-TE policies (aggregation-to-aggregation). The SR-TE policies are nested within another layer of policies using the BSIDs, resulting in seamless end-to-end SR-TE policies.
- Label stack compression—If the label-stack size required for an SR-TE policy exceeds the platform capability, the SR-TE policy can be seamlessly stitched to, or nested within, other SR-TE policies using a binding segment.
- BGP SR-TE Dynamic—The head-end steers the packet into a BGP-based FIB entry whose next hop is a binding-SID.

Stitching SR-TE Polices Using Binding SID: Example

In this intra-domain example, three SR-TE policies are stitched together to form a seamless end-to-end path from node 1 to node 10.

Figure 3: Intra-Domain Topology



Step 1 Configure an SR-TE policy on node 5 to node 10 via node 9. Node 5 automatically allocates a binding-SID (24012) for the SR-TE policy.

Example:

```
RP/0/0/CPU0:xrvr-5(config)# explicit-path name PATH5-9_10
RP/0/0/CPU0:xrvr-5(config-expl-path)# index 10 next-address strict ipv4 unicast 192.168.59.9
RP/0/0/CPU0:xrvr-5(config-expl-path)# index 20 next-address strict ipv4 unicast 10.1.1.10
RP/0/0/CPU0:xrvr-5(config-expl-path)# exit

RP/0/0/CPU0:xrvr-5(config)# interface tunnel-te1
RP/0/0/CPU0:xrvr-5(config-if)# ipv4 unnumbered Loopback0
RP/0/0/CPU0:xrvr-5(config-if)# destination 10.1.1.10
RP/0/0/CPU0:xrvr-5(config-if)# path-option 1 explicit name PATH5-9_10 segment-routing
```

```
RP/0/0/CPU0:xrvr-5(config-if)# commit

RP/0/0/CPU0:xrvr-5# show mpls traffic-eng tunnels 1 detail
Name: tunnel-te1  Destination: 10.1.1.10  Ifhandle:0x680
  Signalled-Name: xrvr-5_t1
  Status:
    Admin:    up Oper:   up   Path:  valid   Signalling: connected
    path option 1, (Segment-Routing) type dynamic (Basis for Setup, path weight 10)
<...>
 Binding SID: 24012
<...>
    Segment-Routing Path Info (IS-IS 1 level-2)
      Segment0[Link]: 192.168.59.5 - 192.168.59.9, Label: 24007
      Segment1[Node]: 10.1.1.10, Label: 16010
```

Step 2 Configure an SR-TE policy on node 3 to node 5 via node 4 and Link4-6, and push the binding-SID of the SR-TE policy at node 5 (24012) to stitch to the SR-TE policy on node 5. Node 3 automatically allocates a binding-SID (24008) for this SR-TE policy.

Example:

```
RP/0/0/CPU0:xrvr-3(config)# explicit-path name PATH4_4-6_5_BSID
RP/0/0/CPU0:xrvr-3(config-expl-path)# index 10 next-address strict ipv4 unicast 10.1.1.4
RP/0/0/CPU0:xrvr-3(config-expl-path)# index 20 next-address strict ipv4 unicast 192.168.46.6
RP/0/0/CPU0:xrvr-3(config-expl-path)# index 30 next-address strict ipv4 unicast 10.1.1.5
RP/0/0/CPU0:xrvr-3(config-expl-path)# index 40 next-label 24012
RP/0/0/CPU0:xrvr-3(config-expl-path)# exit

RP/0/0/CPU0:xrvr-3(config)# interface tunnel-te1
RP/0/0/CPU0:xrvr-3(config-if)# ipv4 unnumbered Loopback0
RP/0/0/CPU0:xrvr-3(config-if)# destination 10.1.1.10
RP/0/0/CPU0:xrvr-3(config-if)# path-option 1 explicit name PATH4_4-6_5_BSID segment-routing
RP/0/0/CPU0:xrvr-3(config-if)# commit

RP/0/0/CPU0:xrvr-3# show mpls traffic-eng tunnels 1 detail
Name: tunnel-te1  Destination: 10.1.1.10  Ifhandle:0x780
  Signalled-Name: xrvr-3_t1
  Status:
    Admin:    up Oper:   up   Path:  valid   Signalling: connected
    path option 1, (Segment-Routing) type explicit PATH4_6_5 (Basis for Setup)
<...>
 Binding SID: 24008
<...>
    Segment-Routing Path Info (IS-IS 1 level-2)
      Segment0[Node]: 10.1.1.4, Label: 16004
      Segment1[Link]: 192.168.46.4 - 192.168.46.6, Label: 24003
      Segment2[Node]: 10.1.1.5, Label: 16005
      Segment3[ - ]: Label: 24012
```

Step 3 Configure an SR-TE policy on node 1 to node 3 and push the binding-SID of the SR-TE policy at node 3 (24008) to stitch to the SR-TE policy on node 3.

Example:

```
RP/0/0/CPU0:xrvr-1(config)# explicit-path name PATH3_BSID
RP/0/0/CPU0:xrvr-1(config-expl-path)# index 10 next-address strict ipv4 unicast 10.1.1.3
RP/0/0/CPU0:xrvr-1(config-expl-path)# index 20 next-label 24008
RP/0/0/CPU0:xrvr-1(config-expl-path)# exit

RP/0/0/CPU0:xrvr-1(config)# interface tunnel-te1
```

```
RP/0/0/CPU0:xrvr-1(config-if)# ipv4 unnumbered Loopback0
RP/0/0/CPU0:xrvr-1(config-if)# destination 10.1.1.10
RP/0/0/CPU0:xrvr-1(config-if)# path-option 1 explicit name PATH3_BSID segment-routing
RP/0/0/CPU0:xrvr-1(config-if)# commit

RP/0/0/CPU0:xrvr-1# show mpls traffic-eng tunnels 1 detail
Name: tunnel-te1  Destination: 10.1.1.10  Ifhandle:0x2f80
  Signalled-Name: xrvr-1_t1
  Status:
    Admin:    up Oper:   up   Path:  valid   Signalling: connected
    path option 1, (Segment-Routing) type explicit PATH3_BSID (Basis for Setup)
<...>
 Binding SID: 24002
<...>
    Segment-Routing Path Info (IS-IS 1 level-2)
      Segment0[Node]: 10.1.1.3, Label: 16003
      Segment1[ - ]: Label: 24008
```

The path is a chain of SR-TE policies stitched together using the binding-SIDs, providing a seamless end-to-end path.

```
RP/0/0/CPU0:xrvr-1# traceroute 10.1.1.10
Type escape sequence to abort.
Tracing the route to 10.1.1.10
 1  99.1.2.2 [MPLS: Labels 16003/24008 Exp 0] 29 msec  19 msec  19 msec
 2  99.2.3.3 [MPLS: Label 24008 Exp 0] 29 msec  19 msec  19 msec
 3  99.3.4.4 [MPLS: Labels 24003/16005/24012 Exp 0] 29 msec  19 msec  19 msec
 4  99.4.6.6 [MPLS: Labels 16005/24012 Exp 0] 29 msec  29 msec  19 msec
 5  99.5.6.5 [MPLS: Label 24012 Exp 0] 29 msec  29 msec  19 msec
 6  99.5.9.9 [MPLS: Label 16010 Exp 0] 19 msec  19 msec  19 msec
 7  99.9.10.10 29 msec  19 msec  19 msec
```

CHAPTER 8

Configure Segment Routing Path Computation Element

The Segment Routing Path Computation Element (SR-PCE) provides stateful PCE functionality by extending the existing IOS-XR PCEP functionality with additional capabilities. SR-PCE is supported on the MPLS data plane and IPv4 control plane.

Note The Cisco IOS XRv 9000 is the recommended platform to act as the SR-PCE. Refer to the Cisco IOS XRv 9000 Router Installation and Configuration Guide for more information.

About SR-PCE

The path computation element protocol (PCEP) describes a set of procedures by which a path computation client (PCC) can report and delegate control of head-end label switched paths (LSPs) sourced from the PCC to a PCE peer. The PCE can request the PCC to update and modify parameters of LSPs it controls. The stateful model also enables a PCC to allow the PCE to initiate computations allowing the PCE to perform network-wide orchestration.

SR-PCE learns topology information by way of IGP (OSPF or IS-IS) or through BGP Link-State (BGP-LS).

SR-PCE is capable of computing paths using the following methods:

- TE metric—SR-PCE uses the TE metric in its path calculations to optimize cumulative TE metric.
- IGP metric—SR-PCE uses the IGP metric in its path calculations to optimize reachability.
- LSP Disjointness—SR-PCE uses the path computation algorithms to compute a pair of disjoint LSPs. The disjoint paths can originate from the same head-end or different head-ends. Disjoint level refers to the type of resources that should not be shared by the two computed paths. SR-PCE supports the following disjoint path computations:
    - Link – Specifies that links are not shared on the computed paths.
    - Node – Specifies that nodes are not shared on the computed paths.

- SRLG – Specifies that links with the same SRLG value are not shared on the computed paths.
- SRLG-node – Specifies that SRLG and nodes are not shared on the computed paths.

When the first request is received with a given disjoint-group ID, the first LSP is computed, encoding the shortest path from the first source to the first destination. When the second LSP request is received with the same disjoint-group ID, information received in both requests is used to compute two disjoint paths: one path from the first source to the first destination, and another path from the second source to the second destination. Both paths are computed at the same time.

Configure SR-PCE

This task explains how to configure SR-PCE.

Before you begin

The Cisco IOS XRv 9000 is the recommended platform to act as the SR-PCE.

SUMMARY STEPS

1. **configure**
2. **pce**
3. **address ipv4** *address*
4. **state-sync ipv4** *address*
5. **tcp-buffer size** *size*
6. **password** {**clear** | **encrypted**} *password*
7. **segment-routing** {**strict-sid-only** | **te-latency**}
8. **timers**
9. **keepalive** *time*
10. **minimum-peer-keepalive** *time*
11. **reoptimization** *time*
12. **exit**

DETAILED STEPS

	Command or Action	Purpose
Step 1	**configure** **Example:** `RP/0/RP0/CPU0:router# configure`	Enters mode.
Step 2	**pce** **Example:** `RP/0/RP0/CPU0:router(config)# pce`	Enables PCE and enters PCE configuration mode.

	Command or Action	Purpose
Step 3	**address ipv4** *address* **Example:** `RP/0/RP0/CPU0:router(config-pce)# address ipv4 192.168.0.1`	Configures a PCE IPv4 address.
Step 4	**state-sync ipv4** *address* **Example:** `RP/0/RP0/CPU0:router(config-pce)# state-sync ipv4 192.168.0.3`	Configures the remote peer for state synchronization.
Step 5	**tcp-buffer size** *size* **Example:** `RP/0/RP0/CPU0:router(config-pce)# tcp-buffer size 1024000`	Configures the transmit and receive TCP buffer size for each PCEP session, in bytes. The default buffer size is 256000. The valid range is from 204800 to 1024000.
Step 6	**password** {**clear** | **encrypted**} *password* **Example:** `RP/0/RP0/CPU0:router(config-pce)# password encrypted pwd1`	Enables TCP MD5 authentication for all PCEP peers. Any TCP segment coming from the PCC that does not contain a MAC matching the configured password will be rejected. Specify if the password is encrypted or clear text.
Step 7	**segment-routing** {**strict-sid-only** | **te-latency**} **Example:** `RP/0/RP0/CPU0:router(config-pce)# segment-routing strict-sid-only`	Configures the segment routing algorithm to use strict SID or TE latency. **Note** This setting is global and applies to all LSPs that request a path from this controller.
Step 8	**timers** **Example:** `RP/0/RP0/CPU0:router(config-pce)# timers`	Enters timer configuration mode.
Step 9	**keepalive** *time* **Example:** `RP/0/RP0/CPU0:router(config-pce-timers)# keepalive 60`	Configures the timer value for locally generated keep-alive messages. The default time is 30 seconds.

	Command or Action	Purpose
Step 10	**minimum-peer-keepalive** *time* **Example:** `RP/0/RP0/CPU0:router(config-pce-timers)# minimum-peer-keepalive 30`	Configures the minimum acceptable keep-alive timer that the remote peer may propose in the PCEP OPEN message during session establishment. The default time is 20 seconds.
Step 11	**reoptimization** *time* **Example:** `RP/0/RP0/CPU0:router(config-pce-timers)# reoptimization 600`	Configures the re-optimization timer. The default timer is 1800 seconds.
Step 12	**exit** **Example:** `RP/0/RP0/CPU0:router(config-pce-timers)# exit`	Exits timer configuration mode and returns to PCE configuration mode.

Configure the Disjoint Policy (Optional)

This task explains how to configure the SR-PCE to compute disjointness for a pair of LSPs signaled by PCCs that do not include the PCEP association group-ID object in their PCEP request. This can be beneficial for deployments where PCCs do not support this PCEP object or when the network operator prefers to manage the LSP disjoint configuration centrally.

SUMMARY STEPS

1. **disjoint-path**
2. **group-id** *value* **type** {**link** | **node** | **srlg** | **srlg-node**} [**sub-id** *value*]
3. **strict**
4. **lsp** {**1** | **2**} **pcc ipv4** *address* **lsp-name** *lsp_name* [**shortest-path**]

DETAILED STEPS

	Command or Action	Purpose
Step 1	**disjoint-path** **Example:** `RP/0/RP0/CPU0:router(config-pce)# disjoint-path`	Enters disjoint configuration mode.
Step 2	**group-id** *value* **type** {**link** | **node** | **srlg** | **srlg-node**} [**sub-id** *value*] **Example:**	Configures the disjoint group ID and defines the preferred level of disjointness (the type of resources that should not be shared by the two paths):

	Command or Action	Purpose
	`RP/0/RP0/CPU0:router(config-pce-disjoint)# group-id 1 type node sub-id 1`	- **link**—Specifies that links are not shared on the computed paths. - **node**—Specifies that nodes are not shared on the computed paths. - **srlg**—Specifies that links with the same SRLG value are not shared on the computed paths. - **srlg-node**—Specifies that SRLG and nodes are not shared on the computed paths. If a pair of paths that meet the requested disjointness level cannot be found, then the paths will automatically fallback to a lower level: - If the requested disjointness level is SRLG or node, then link-disjoint paths will be computed. - If the requested disjointness level was link, or if the first fallback from SRLG or node disjointness failed, then the lists of segments encoding two shortest paths, without any disjointness constraint, will be computed.
Step 3	**strict** **Example:** `RP/0/RP0/CPU0:router(config-pce-disjoint)# strict`	(Optional) Prevents the automatic fallback behavior of the preferred level of disjointness. If a pair of paths that meet the requested disjointness level cannot be found, the disjoint calculation terminates and no new path is provided. The existing path is not modified.

Step 4	**lsp** {**1** | **2**} **pcc ipv4** *address* **lsp-name** *lsp_name* [**shortest-path**]
	Example:
	`RP/0/RP0/CPU0:router(config-pce-disjoint)# lsp 1 pcc ipv4 192.168.0.1 lsp-name rtrA_t1 shortest-path`
	`RP/0/RP0/CPU0:router(config-pce-disjoint)# lsp 2 pcc ipv4 192.168.0.5 lsp-name rtrE_t2`
		Adds LSPs to the disjoint group.
		The **shortest-path** keyword forces one of the disjoint paths to follow the shortest path from the source to the destination. This option can only be applied to the the first LSP specified.

CHAPTER 9

Configure Topology-Independent Loop-Free Alternate (TI-LFA)

Topology-Independent Loop-Free Alternate (TI-LFA) uses segment routing to provide link protection in topologies where other fast reroute techniques cannot provide protection.

- Classic Loop-Free Alternate (LFA) is topology dependent, and therefore cannot protect all destinations in all networks. A limitation of LFA is that, even if one or more LFAs exist, the optimal LFA may not always be provided.
- Remote LFA (RLFA) extends the coverage to 90-95% of the destinations, but it also does not always provide the most desired repair path. RLFA also adds more operational complexity by requiring a targeted LDP session to the RLFAs to protect LDP traffic.

TI-LFA provides a solution to these limitations while maintaining the simplicity of the IPFRR solution.

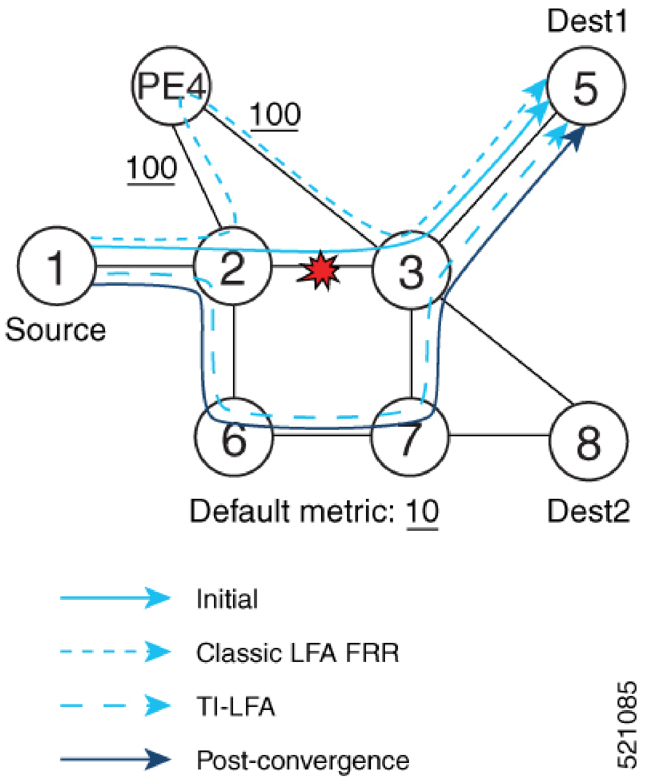
The goal of TI-LFA is to reduce the packet loss that results while routers converge after a topology change due to a link failure. Rapid failure repair (< 50 msec) is achieved through the use of pre-calculated backup paths that are loop-free and safe to use until the distributed network convergence process is completed.

The optimal repair path is the path that the traffic will eventually follow after the IGP has converged. This is called the post-convergence path. This path is preferred for the following reasons:

- Optimal for capacity planning — During the capacity-planning phase of the network, the capacity of a link is provisioned while taking into consideration that such link with be used when other links fail.
- Simple to operate — There is no need to perform a case-by-case adjustments to select the best LFA among multiple candidate LFAs.
- Fewer traffic transitions — Since the repair path is equal to the post-convergence path, the traffic switches paths only once.

The following topology illustrates the optimal and automatic selection of the TI-LFA repair path.

Figure 4: TI-LFA Repair Path



Node 2 protects traffic to destination Node 5.

With classic LFA, traffic would be steered to Node 4 after a failure of the protected link. This path is not optimal, since traffic is routed over edge node Node 4 that is connected to lower capacity links.

TI-LFA calculates a post-convergence path and derives the segment list required to steer packets along the post-convergence path without looping back.

In this example, if the protected link fails, the shortest path from Node2 to Node5 would be:

Node2 → Node6 → Node7 → Node3 → Node5

Node7 is the PQ-node for destination Node5. TI-LFA encodes a single segment (prefix SID of Node7) in the header of the packets on the repair path.

- Behaviors and Limitations of TI-LFA, on page 56
- Configuring TI-LFA for IS-IS, on page 57
- Configuring TI-LFA for OSPF, on page 58

Behaviors and Limitations of TI-LFA

The behaviors and limitations of TI-LFA are listed below:

- TI-LFA protects unlabeled IPv4 traffic.
- TI-LFA does not protect unlabeled IPv6 traffic.

Configuring TI-LFA for IS-IS

This task describes how to enable per-prefix Topology Independent Loop-Free Alternate (TI-LFA) computation to converge traffic flows around link failures.

Before you begin

Ensure that the following topology requirements are met:

- Router interfaces are configured as per the topology.
- Routers are configured with IS-IS.
- Segment routing for IS-IS is configured. See Enabling Segment Routing for IS-IS Protocol, on page 11.
- Enter the following commands in global configuration mode:

```
Router(config)# ipv4 unnumbered mpls traffic-eng Loopback0
Router(config)# mpls traffic-eng
Router(config-mpls-te)# exit
Router(config)#
```

SUMMARY STEPS

1. **configure**
2. **router isis** *instance-id*
3. **interface** *type interface-path-id*
4. **address-family ipv4** [**unicast**]
5. **fast-reroute per-prefix**
6. **fast-reroute per-prefix ti-lfa**

DETAILED STEPS

	Command or Action	Purpose
Step 1	**configure** **Example:** `RP/0/RP0/CPU0:router# configure`	Enters mode.
Step 2	**router isis** *instance-id* **Example:** `RP/0/RP0/CPU0:router(config)# router isis 1`	Enables IS-IS routing for the specified routing instance, and places the router in router configuration mode. **Note** You can change the level of routing to be performed by a particular routing instance by using the **is-type** router configuration command.
Step 3	**interface** *type interface-path-id* **Example:** `RP/0/RP0/CPU0:router(config-isis)# interface`	Enters interface configuration mode.

	Command or Action	Purpose
	`GigabitEthernet0/0/2/1`	
Step 4	**address-family ipv4** [**unicast**] **Example:** `RP/0/RP0/CPU0:router(config-isis-if)#` **`address-family ipv4 unicast`**	Specifies the IPv4 address family, and enters router address family configuration mode.
Step 5	**fast-reroute per-prefix** **Example:** `RP/0/RP0/CPU0:router(config-isis-if-af)#` **`fast-reroute per-prefix`**	Enables per-prefix fast reroute.
Step 6	**fast-reroute per-prefix ti-lfa** **Example:** `RP/0/RP0/CPU0:router(config-isis-if-af)#` **`fast-reroute per-prefix ti-lfa`**	Enables per-prefix TI-LFA fast reroute link protection.

TI-LFA has been successfully configured for segment routing.

Configuring TI-LFA for OSPF

This task describes how to enable per-prefix Topology Independent Loop-Free Alternate (TI-LFA) computation to converge traffic flows around link failures.

Note TI-LFA can be configured on the instance, area, or interface. When configured on the instance or area, all interfaces in the instance or area inherit the configuration.

Before you begin

Ensure that the following topology requirements are met:

- Router interfaces are configured as per the topology.
- Routers are configured with OSPF.
- Segment routing for OSPF is configured. See Enabling Segment Routing for OSPF Protocol, on page 19.
- Enter the following commands in global configuration mode:

```
Router(config)# ipv4 unnumbered mpls traffic-eng Loopback0
Router(config)# mpls traffic-eng
```

```
Router(config-mpls-te)# exit
Router(config)#
```

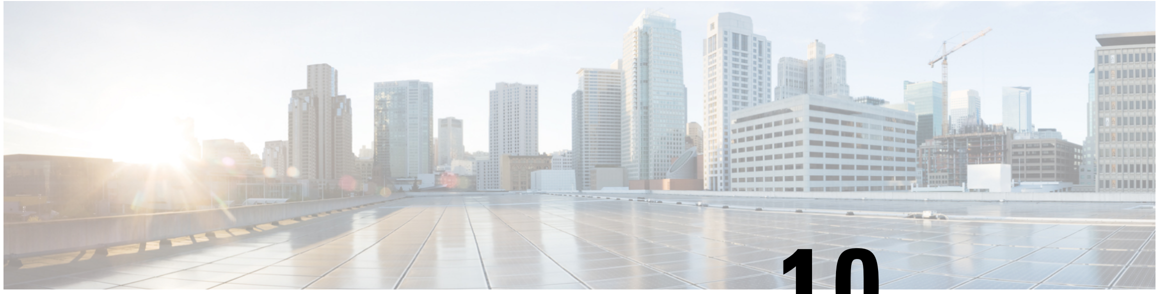
SUMMARY STEPS

1. **configure**
2. **router ospf** *process-name*
3. **area** *area-id*
4. **interface** *type interface-path-id*
5. **fast-reroute per-prefix**
6. **fast-reroute per-prefix ti-lfa**

DETAILED STEPS

	Command or Action	Purpose
Step 1	**configure** **Example:** `RP/0/RP0/CPU0:router# configure`	Enters mode.
Step 2	**router ospf** *process-name* **Example:** `RP/0/RP0/CPU0:router(config)# router ospf 1`	Enables OSPF routing for the specified routing process, and places the router in router configuration mode.
Step 3	**area** *area-id* **Example:** `RP/0/RP0/CPU0:router(config-ospf)# area 1`	Enters area configuration mode.
Step 4	**interface** *type interface-path-id* **Example:** `RP/0/RP0/CPU0:router(config-ospf-ar)# interface GigabitEthernet0/0/2/1`	Enters interface configuration mode.
Step 5	**fast-reroute per-prefix** **Example:** `RP/0/RP0/CPU0:router(config-ospf-ar-if)# fast-reroute per-prefix`	Enables per-prefix fast reroute.
Step 6	**fast-reroute per-prefix ti-lfa** **Example:** `RP/0/RP0/CPU0:router(config-ospf-ar-if)# fast-reroute per-prefix ti-lfa`	Enables per-prefix TI-LFA fast reroute link protection.

TI-LFA has been successfully configured for segment routing.



CHAPTER 10

Using Segment Routing OAM

Segment Routing Operations, Administration, and Maintenance (OAM) helps service providers to monitor label-switched paths (LSPs) and quickly isolate forwarding problems to assist with fault detection and troubleshooting in the network. The Segment Routing OAM feature provides support for Nil-FEC (forwarding equivalence classes) LSP Ping and Traceroute functionality.

MPLS LSP Ping and Traceroute Nil FEC Target

The Nil-FEC LSP ping and traceroute operations are extensions of regular MPLS ping and traceroute.

Nil-FEC LSP Ping/Traceroute functionality supports segment routing and MPLS Static. It also acts as an additional diagnostic tool for all other LSP types. This feature allows operators to provide the ability to freely test any label stack by allowing them to specify the following:

- label stack
- outgoing interface
- nexthop address

In the case of segment routing, each segment nodal label and adjacency label along the routing path is put into the label stack of an echo request message from the initiator Label Switch Router (LSR); MPLS data plane forwards this packet to the label stack target, and the label stack target sends the echo message back.

The following table shows the syntax for the ping and traceroute commands.

Table 2: LSP Ping and Traceroute Nil FEC Commands

Command Syntax
ping mpls nil-fec labels {*label*[,*label*]} [**output** {**interface** *tx-interface*} [**nexthop** *nexthop-ip-addr*]]
traceroute mpls nil-fec labels {*label*[,*label*]} [**output** {**interface** *tx-interface*} [**nexthop** *nexthop-ip-addr*]]

Examples: LSP Ping and Traceroute for Nil_FEC Target

These examples use the following topology:

```
Node loopback IP address: 172.18.1.3   172.18.1.4   172.18.1.5   172.18.1.7
Node label:                            16004        16005        16007
Nodes:                    Arizona ---- Utah ------- Wyoming ---- Texas

Interface:            GigabitEthernet0/2/0/1   GigabitEthernet0/2/0/1
Interface IP address:         10.1.1.3              10.1.1.4
```

```
RP/0/RP0/CPU0:router-utah# show mpls forwarding

Tue Jul  5 13:44:31.999 EDT
Local  Outgoing    Prefix             Outgoing     Next Hop        Bytes
Label  Label       or ID              Interface                    Switched
------ ----------- ------------------ ------------ --------------- ------------
16004  Pop         No ID              Gi0/2/0/1    10.1.1.4        1392
       Pop         No ID              Gi0/2/0/2    10.1.2.2        0
16005  16005       No ID              Gi0/2/0/0    10.1.1.4        0
       16005       No ID              Gi0/2/0/1    10.1.2.2        0
16007  16007       No ID              Gi0/2/0/0    10.1.1.4        4752
       16007       No ID              Gi0/2/0/1    10.1.2.2        0
24000  Pop         SR Adj (idx 0)     Gi0/2/0/0    10.1.1.4        0
24001  Pop         SR Adj (idx 2)     Gi0/2/0/0    10.1.1.4        0
24002  Pop         SR Adj (idx 0)     Gi0/2/0/1    10.1.2.2        0
24003  Pop         SR Adj (idx 2)     Gi0/2/0/1    10.1.2.2        0
24004  Pop         No ID              tt10         point2point     0
24005  Pop         No ID              tt11         point2point     0
24006  Pop         No ID              tt12         point2point     0
24007  Pop         No ID              tt13         point2point     0
24008  Pop         No ID              tt30         point2point     0
```

Ping Nil FEC Target

```
RP/0/RP0/CPU0:router-arizona# ping mpls nil-fec labels 16005,16007 output interface
GigabitEthernet 0/2/0/1 nexthop 10.1.1.4 repeat 1
Sending 1, 72-byte MPLS Echos with Nil FEC labels 16005,16007,
     timeout is 2 seconds, send interval is 0 msec:

Codes: '!' - success, 'Q' - request not sent, '.' - timeout,
  'L' - labeled output interface, 'B' - unlabeled output interface,
  'D' - DS Map mismatch, 'F' - no FEC mapping, 'f' - FEC mismatch,
  'M' - malformed request, 'm' - unsupported tlvs, 'N' - no label entry,
  'P' - no rx intf label prot, 'p' - premature termination of LSP,
  'R' - transit router, 'I' - unknown upstream index,
  'l' - Label switched with FEC change, 'd' - see DDMAP for return code,
  'X' - unknown return code, 'x' - return code 0

Type escape sequence to abort.
!
Success rate is 100 percent (1/1), round-trip min/avg/max = 1/1/1 ms
 Total Time Elapsed 0 ms
```

Traceroute Nil FEC Target

```
RP/0/RP0/CPU0:router-arizona# traceroute mpls nil-fec labels 16005,16007 output interface
GigabitEthernet 0/2/0/1 nexthop 10.1.1.4
Tracing MPLS Label Switched Path with Nil FEC labels 16005,16007, timeout is 2 seconds

Codes: '!' - success, 'Q' - request not sent, '.' - timeout,
  'L' - labeled output interface, 'B' - unlabeled output interface,
  'D' - DS Map mismatch, 'F' - no FEC mapping, 'f' - FEC mismatch,
  'M' - malformed request, 'm' - unsupported tlvs, 'N' - no label entry,
  'P' - no rx intf label prot, 'p' - premature termination of LSP,
  'R' - transit router, 'I' - unknown upstream index,
  'l' - Label switched with FEC change, 'd' - see DDMAP for return code,
  'X' - unknown return code, 'x' - return code 0

Type escape sequence to abort.
  0 10.1.1.3 MRU 1500 [Labels: 16005/16007/explicit-null Exp: 0/0/0]
L 1 10.1.1.4 MRU 1500 [Labels: implicit-null/16007/explicit-null Exp: 0/0/0] 1 ms
L 2 10.1.1.5 MRU 1500 [Labels: implicit-null/explicit-null Exp: 0/0] 1 ms
! 3 10.1.1.7 1 ms
```